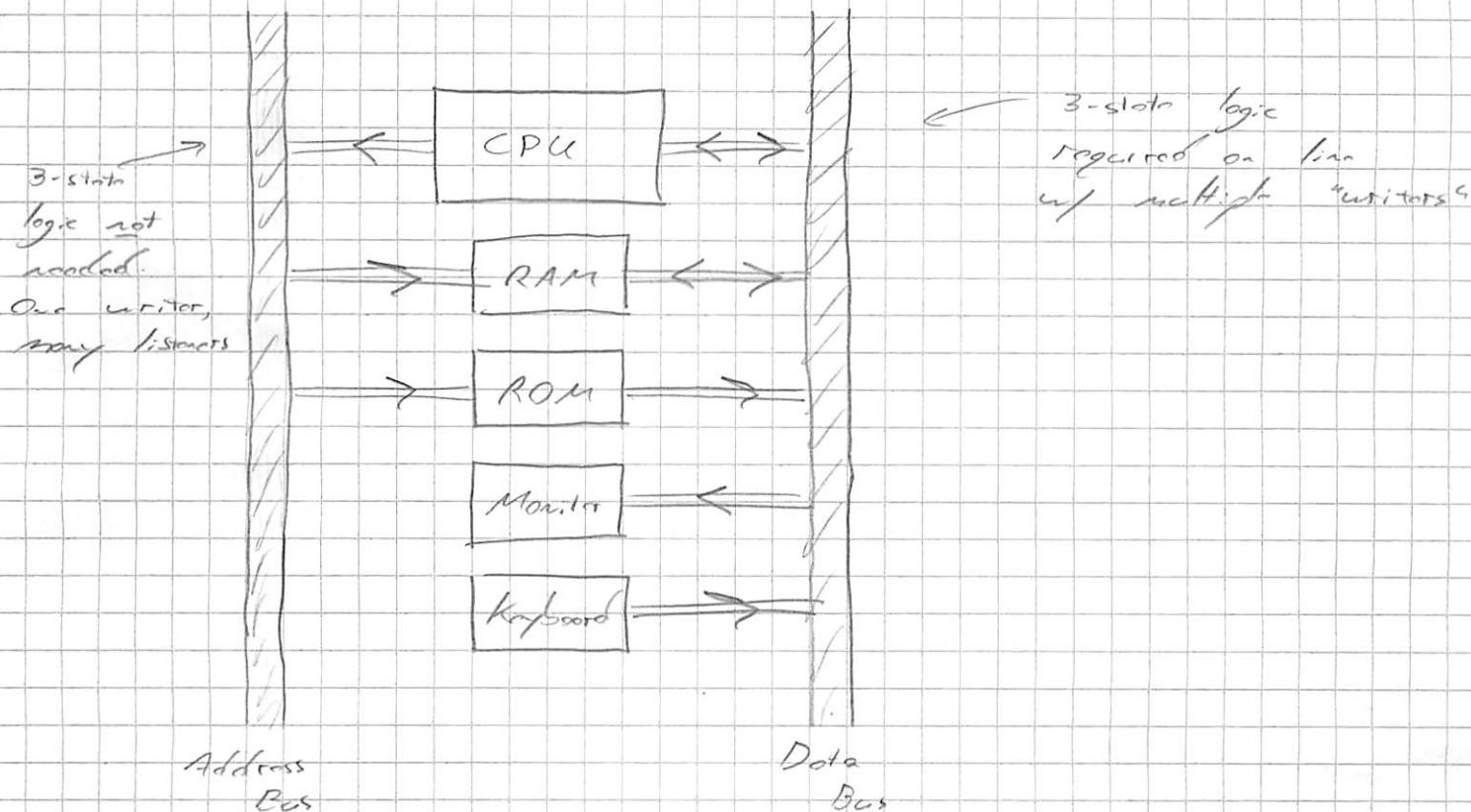


Digital III & Digital meets Analog

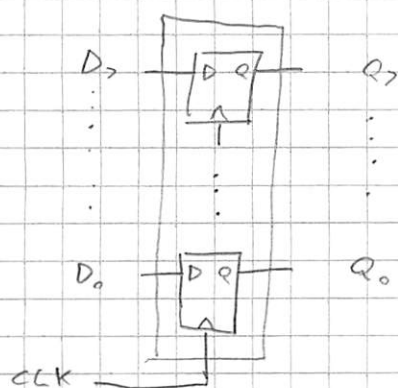
Computer Organization :



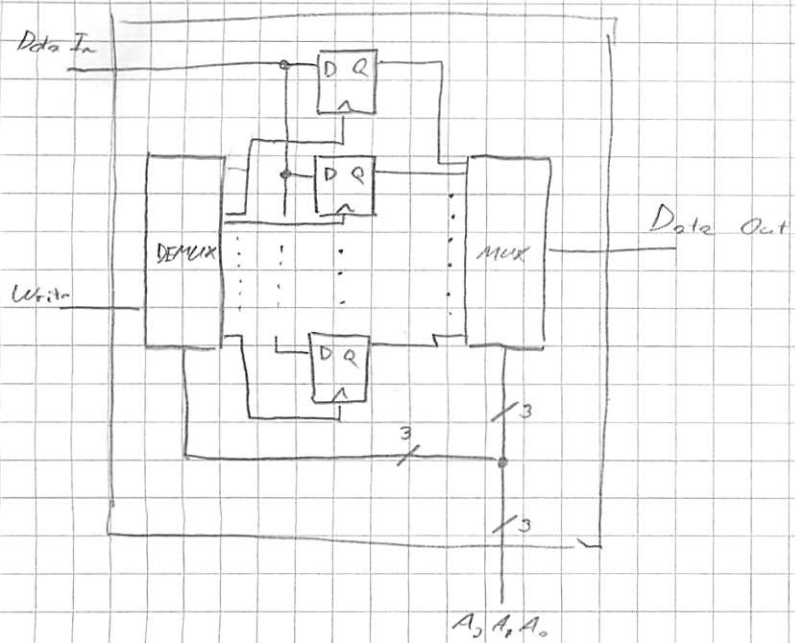
Memory

Array of Flip-flops (or other device which can hold bits of information). Strong information

8-bit memory : register (collection of flip-flops)



	<u>Pins</u>
data in	8
data out	8
clock/write	1
power/gnd	2
	19

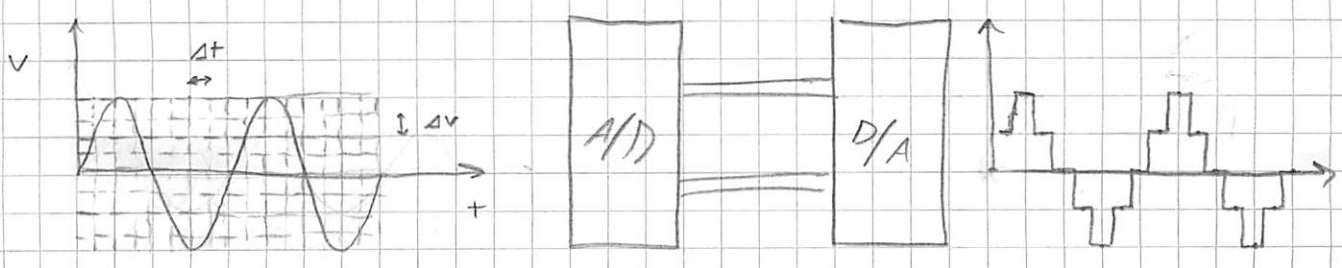


	Pins
data in	1
data out	1
clock/pwr	1
power/gnd	2
address	3
	8

fewer pins ... as memory becomes larger & larger the 2-d architecture becomes more and more favorable.

D-to-A & A-to-D
 > DAC, D/A > ADC, A/D

D-to-A (Digital to Analog Conversion)



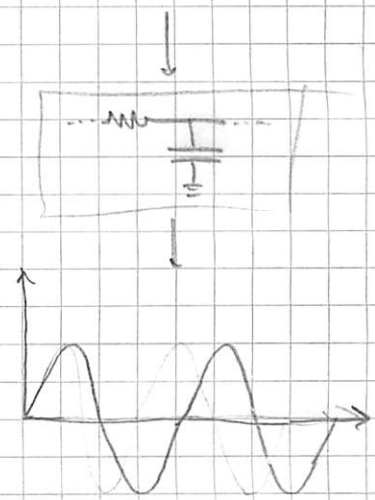
$\Delta V \rightarrow$ voltage slice

$V \leftrightarrow$ binary #

8-bit encoding : $2^8 = 256$ slices

$$\Delta V = \frac{V_{Fullscale}}{256}$$

$\frac{\Delta V}{2} =$ voltage resolution



How many bits required for 0.01% resolution? ③

$$0.01\% = \frac{1}{10000}$$

→ For 5V full scale,

$$\rightarrow 0.5 \text{ mV}$$

$$\begin{array}{l} 9.9115 \rightarrow \\ 9.9105 \rightarrow \end{array} \quad \begin{array}{l} 9.912 \text{ V} \\ 9.911 \text{ V} \\ 9.910 \text{ V} \end{array} \quad \Delta V = 1 \text{ mV}$$

→ corresponds to

$$5000 \text{ slices} \approx 2^{12}$$

∴ 12-bit encoding

$\Delta t \rightarrow$ sampling period

$\frac{1}{\Delta t} \rightarrow$ sampling rate

f_s

Just over 2 samples per period are required to

describe a sine wave — Nyquist

Theorem

⇒ If $x(t)$ contains no frequencies higher than B , it can be completely determined by measuring its value at a series of points spaced by $\Delta t = \frac{1}{2B}$.

$$\rightarrow f_s \geq 2B$$

$$f(t) = \frac{1}{2\pi} \int_{-\infty}^{\infty} F(\omega) e^{i\omega t} d\omega$$

No freqs higher than B :

$$f(t) = \frac{1}{2\pi} \int_{-2\pi B}^{2\pi B} F(\omega) e^{i\omega t} d\omega$$

$$F(\omega) = \sum_{n=-\infty}^{\infty} c_n e^{in\omega}$$

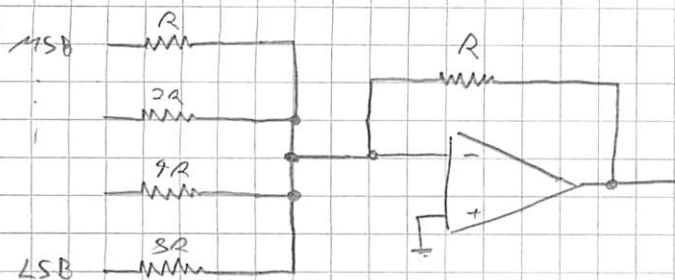
$$c_n = \frac{1}{2\pi} \int_{-\pi}^{\pi} F(\omega) e^{-in\omega} d\omega$$

$$f(t) = \frac{1}{2\pi} \int_{-\infty}^{\infty} F(\omega) e^{i\omega t} d\omega$$

$$f(t) = \frac{1}{2\pi} \int_{-\infty}^{\infty} \sum_{n=-\infty}^{\infty} c_n e^{in\omega} e^{i\omega t} d\omega$$

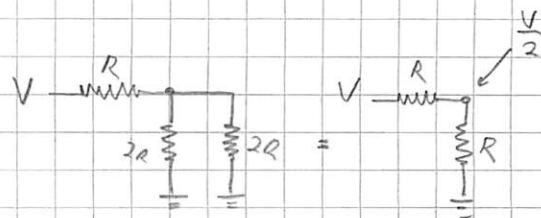
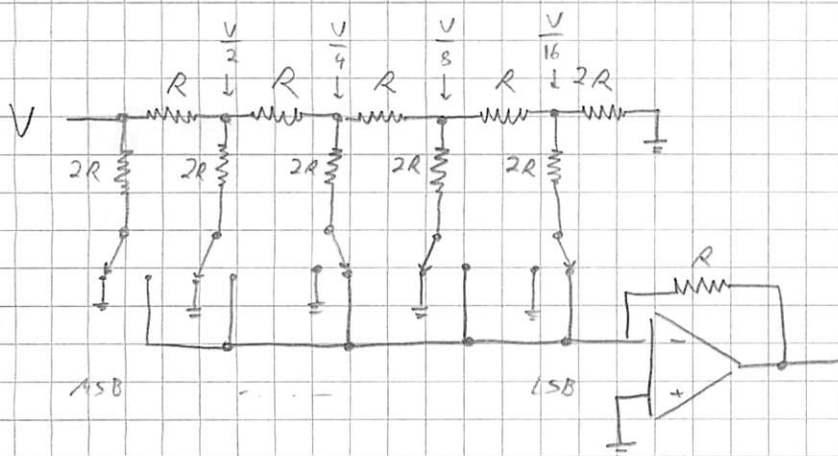
D-to-A

4



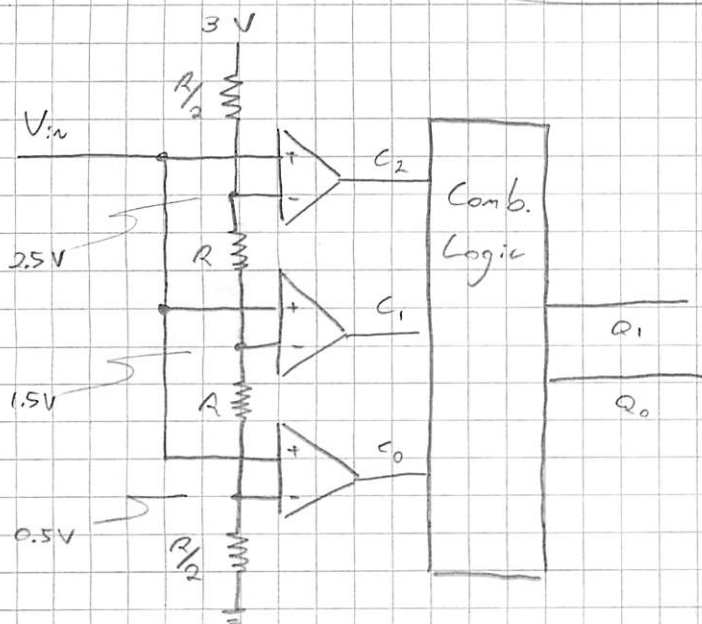
This method becomes intractable for large bit encoding... Resistors w/ ratios of 2000:1 for a 10-bit converter...

Solution: R-2R Ladder



A-to-D

Parallel Encoders (Flash encoder)

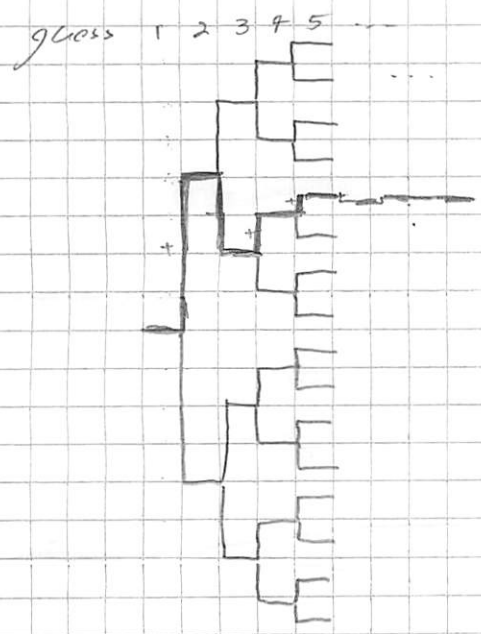
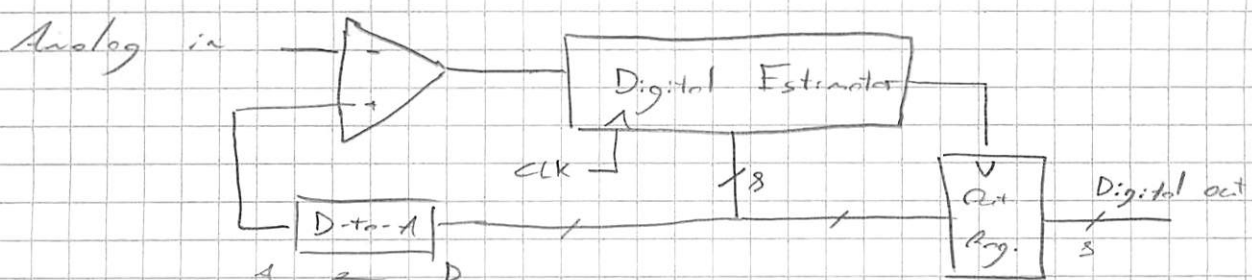
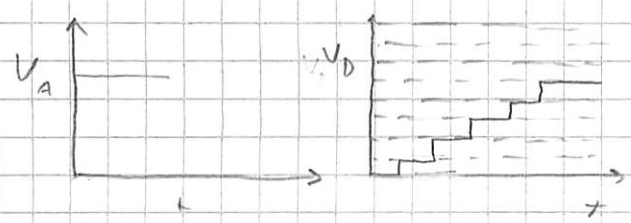
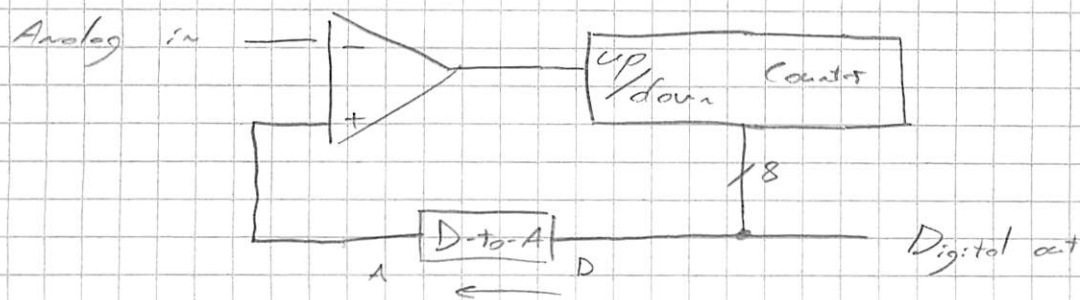


Thresholds	V_{in}	Comp. Output C_2, C_1, C_0	Code Q_1, Q_0
3V	2.8V	1 1 1	1 1
2.5V	2.0V	0 1 1	1 0
1.5V	1.0V	0 0 1	0 1
0.5V	0.25V	0 0 0	0 0
0V			

$$2^2 - 1$$

$$2^3 - 1 = 511$$

Closed-loop : Successive - approximation & tracking



Guess #	D ₇	D ₆	D ₅	D ₄	D ₃	D ₂	D ₁	D ₀	
1	0	1	1	1	1	1	1	1	+
2	1	0	1	1	1	1	1	1	-
3	1	0	0	1	1	1	1	1	+
4	1	0	1	0	1	1	1	1	+
5	1	0	1	1	0	1	1	1	-
6	1	0	1	1	0	0	1	1	+
7	1	0	1	1	0	1	0	1	-
8	1	0	1	1	0	1	0	0	

Fourier Transform

①

$$f(t) = \int_{-\infty}^{\infty} \hat{f}(\nu) e^{2\pi i \nu t} d\nu$$

$$\hat{f}(\nu) = \int_{-\infty}^{\infty} f(t) e^{-2\pi i \nu t} dt$$

But we sample the continuous signal $f(t)$ at discrete time intervals: $0, T, 2T, \dots$ w/ $T = 1/\nu_{\text{sampling}}$ and N is the number of samples. A discrete Fourier transform can then be written:

$$\hat{f}(\nu) = \lim_{\substack{T \rightarrow 0 \\ N \rightarrow \infty}} \sum_{n=0}^{N-1} f(nT) e^{-2\pi i n T \nu} \cdot T$$

↑
cont. FT disc. FT $F(\nu)$

$$F(\nu) = T \sum_{n=0}^{N-1} f(nT) e^{-2\pi i n T \nu}$$

$$F(\nu_{N/2} - \nu') = T \sum_{n=0}^{N-1} f(nT) \cdot e^{-2\pi i n T (\nu_{N/2} - \nu')}$$

$$\text{w/ } \nu_{N/2} = \nu_{\text{sampling}}/2 = \frac{1}{2T}$$

$$F(\nu_{N/2} - \nu') = T \sum_{n=0}^{N-1} f(nT) \underbrace{e^{-\pi i n} e^{2\pi i n T \nu'}}_{e^{i\pi n} e^{2\pi i n T \nu'}}$$

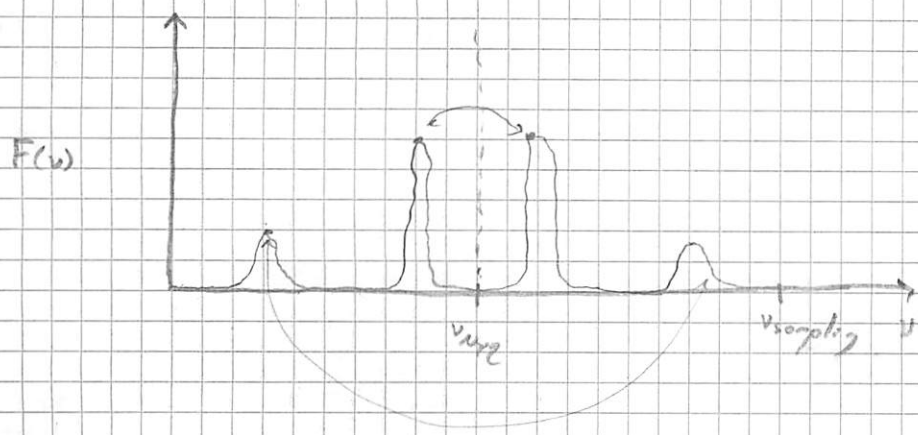
$$e^{i\pi n} e^{2\pi i n T \nu'} = e^{i2\pi (\frac{1}{2T} + \nu') n T} = e^{i2\pi (\nu_{N/2} + \nu') n T}$$

$$F(\nu_{N/2} - \nu') = T \sum_{n=0}^{N-1} f(nT) \cdot e^{2\pi i n T (\nu_{N/2} + \nu')}$$

$$F(\nu_{N/2} - \nu') = \left[T \sum_{n=0}^{N-1} f(nT) e^{-2\pi i n T (\nu_{N/2} + \nu')} \right]^*$$

$$\therefore \boxed{F(\nu_{N/2} - \nu') = F(\nu_{N/2} + \nu')}$$

The previous relation results in so-called Frequency "folding".



Actually the problem is much worse!



$$F(v) = T \sum_{n=0}^{N-1} f(nT) e^{-2\pi i n T v}$$

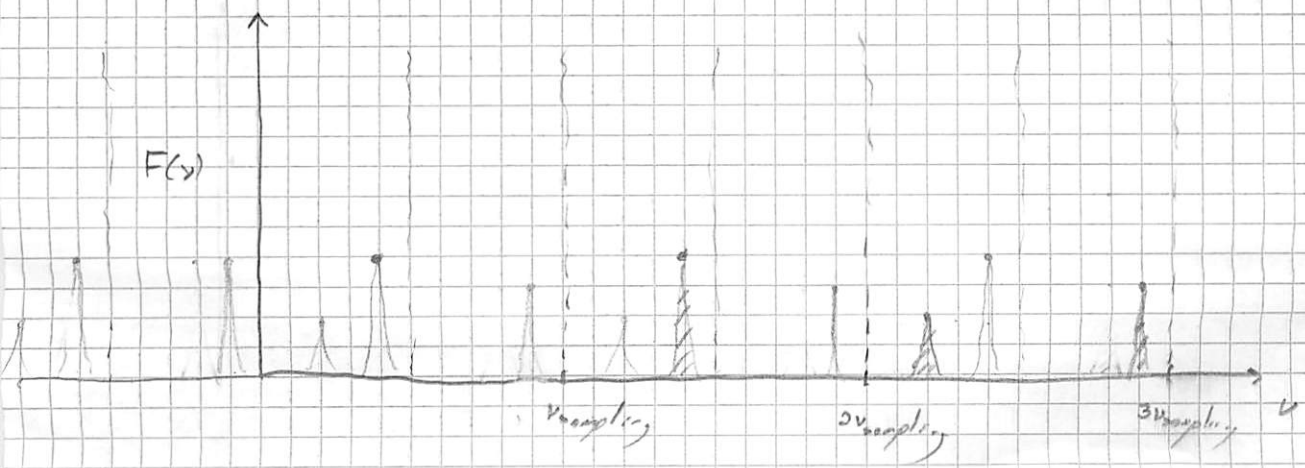
$$F(v + m v_{\text{sampling}}) = T \sum_{n=0}^{N-1} f(nT) e^{-2\pi i n T (v + m v_{\text{sampling}})}$$

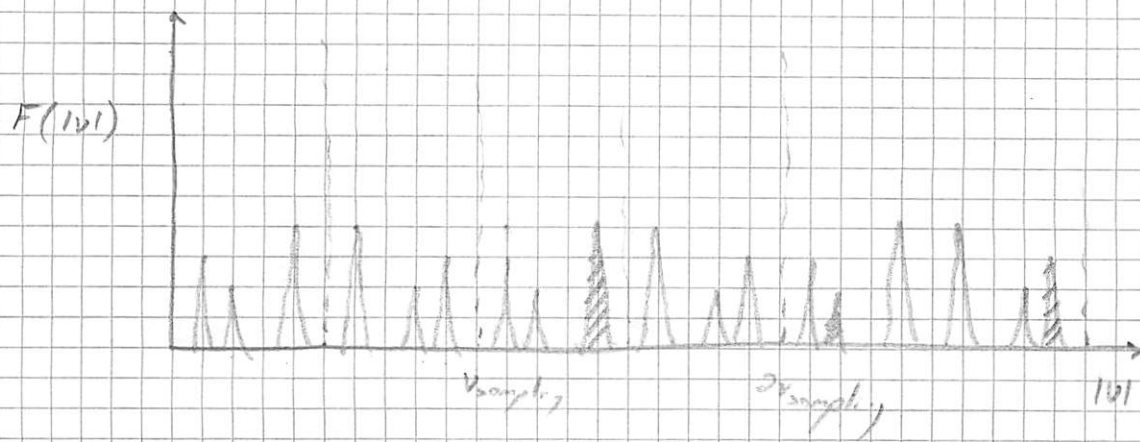
$$e^{-2\pi i n T v} \quad e^{-2\pi i n m}$$

$m = 0, \pm 1, \pm 2, \dots$

$$F(v + m v_{\text{sampling}}) = T \sum_{n=0}^{N-1} f(nT) e^{-2\pi i n T v}$$

$$F(v + m v_{\text{sampling}}) = F(v)$$





What concerns us...

Solution: Filtering for $v \geq v_{\text{sampling}}/2 = v_{\text{Nyq}}$